

5. The *girth* (omkrets in Swedish) of a graph G is the number of edges of the shortest cycle it is possible to find in G and it is denoted $g(G)$. Let, from now on, the graphs G have $e = v + 1$. For such graphs with $v \geq 4$ it is possible to show that

$$g(G) \leq \left\lfloor \frac{2(v+1)}{3} \right\rfloor.$$

The floor function $\lfloor x \rfloor$ gives the largest integer $\leq x$.

- a) Explain why in this type of graphs there must be at least two cycles. Draw a case where the cycles are edge disjoint for $v = 7$ and $e = 8$. (2p)
- b) If the two cycles are edge disjoint what is the maximum girth in a graph with $e = v + 1$? (1p)
- c) Draw a graph with maximal girth for $v = 7$ and $e = 8$. (2p)
6. a) What is the maximal number of edges in a planar graph with v vertices when the regions have degree 3 or more? Use for example $2e = \sum \deg(R_i)$ and Euler's formula. (2p)
- b) Let G be a graph with 11 vertices and let \bar{G} be its complement graph. \bar{G} has the same 11 vertices but the edges from K_{11} that are *not* edges in G . Use the result in (a) above to show that not both G and \bar{G} can be planar. (2p)

Good Luck!